

Context-Free Grammars (CFGs) (1/3)

CFGs offer concise language specifications.

$$S ::= D \mid DS$$

$$D ::= 0 \mid 1$$

A grammar consists of:

- a set of **non-terminals** and a set of **terminals**
- a set of productions (here four):
 - “an S may be replaced by D ”
 - “an S may be replaced by DS ”
 - “a D may be replaced by 0”
 - “a D may be replaced by 1”

□ A (non-terminal) start symbol S , typically written first.

Context-Free Grammars (CFGs) (2/3)

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$$D ::= 0 \mid 1$$

They do so by expressing a set of derivable strings.

Recipe: Repeatedly replace a non-terminal by a right-hand-side until there's only terminals left:

$$S \rightarrow D \rightarrow 0$$

$$S \rightarrow DS \rightarrow 0S \rightarrow 0D \rightarrow 01$$

$$S \rightarrow DS \rightarrow 1S \rightarrow 1DS \rightarrow 10S \rightarrow 10D \rightarrow 101$$

Q: Which language does this grammar represent?

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Q: Which language does this grammar represent?

Exercise: alter the grammar to yield base-10 digits

Context-Free Grammars (CFGs) (3/3)

BNF grammars (short for ‘**Backus-Naur Form**’ after John Backus and Peter Naur) are a particular kind of CFGs.

A CFG can express a ‘**Context-Free Language**’. These are strictly more powerful than, e.g., regular expressions.

Most language specifications today come with a CFG.

Q: Which languages do the following grammars define?

$S ::= 0 \mid 01 \mid 01S$ (terminals: 0, 1)

$S ::= ab \mid aSb$ (terminals: a, b)

$S ::= e \mid S + S \mid S - S$ (terminals: e, +, -)

$S ::= s \mid S ; S \mid \text{if } (e) \{ S \} \mid \text{while } (e) \{ S \}$
(terminals: s, e, ;, if, while, (,), {, })^{3/4}

Other CFG examples

Grammars are used in real-world specifications:

A JSON grammar:

<https://www.json.org/json-en.html>

A grammar from the HTTP 1.1 RFC:

<https://tools.ietf.org/html/rfc2616#page-14>

The Java grammar from the Java Language Specification:

<https://docs.oracle.com/javase/specs/jls/se13/html/jl>